

# A proof of the GM-MDS conjecture

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## Abstract

The GM-MDS conjecture of Dau et al. (ISIT 2014) speculates that the MDS condition, which guarantees the existence of MDS matrices with a prescribed set of zeros over large fields, is in fact sufficient for existence of such matrices over small fields. We prove this conjecture.

## 1 Introduction

Let  $S_1, \dots, S_k \subset [n]$  be a set system with  $k \leq n$ . A natural question, motivated by the study of MDS codes, is when does there exist a  $k \times n$  matrix  $G$  over a finite field  $\mathbb{F}$ , such that all  $k \times k$  minors of  $G$  are full rank, and such that  $G_{i,j} = 0$  whenever  $j \in S_i$ .

If there is no limitation on the field size then the answer is well known. Any row can have at most  $k - 1$  zeros. Any pair of rows at most  $k - 2$  common zeros, and so on. This is the so-called *MDS condition*:

$$\left| \bigcap_{i \in I} S_i \right| \leq k - |I| \quad \forall I \subseteq [k]. \quad (1)$$

Concretely, for any set system which satisfies the MDS condition, and for any field of size  $|\mathbb{F}| > \binom{n-1}{k-1}$ , there exist matrices  $G$  with a prescribed set of zeros given by the set system (see for example [DSDY13] for details). In fact, if  $|\mathbb{F}| \gg \binom{n}{k}$  then one can simply choose a random matrix  $G$  with  $G_{i,j} = 0$  if  $j \in S_i$ , and  $G_{i,j} \in \mathbb{F}$  uniformly for all other entries, and obtain that with high probability all the minors of  $G$  are nonsingular.

A much more subtle question is that is the minimum field size for which such a matrix exists. Dau et al. [DSY14] conjectured that the MDS condition (1) is sufficient for the existence of a matrix  $G$  over much smaller fields.

**Conjecture 1.1** (GM-MDS conjecture [DSY14]). *Let  $S_1, \dots, S_k \subset [n]$  be a set system which satisfies the MDS condition. Then for any field  $\mathbb{F}$  with  $|\mathbb{F}| \geq n + k - 1$ , there exists a  $k \times n$  matrix  $G$  over  $\mathbb{F}$  with  $G_{i,j} = 0$  whenever  $j \in S_i$ , such that all  $k \times k$  minors of  $G$  are nonsingular.*

We prove Conjecture 1.1 in this work. First, we describe an algebraic framework introduced by Dau et al. [DSY14] towards proving Conjecture 1.1.

## 1.1 The algebraic GM-MDS conjecture

Dau et al. [DSY14] formulated an algebraic conjecture that implies Conjecture 1.1: if  $S_1, \dots, S_k$  is a set system that satisfies (1), then there exists a Generalized Reed-Muller code with zeros in locations prescribed by the set system. Otherwise put, there exists  $G = AV$  where  $A$  is an invertible  $k \times k$  matrix and  $V$  a  $k \times n$  Vandermonde matrix, such that  $G_{i,j} = 0$  when  $j \in S_i$ . Before we explain these ideas further, we first set up some notations.

Let  $\mathbb{F}$  be a finite field, and let  $x, a_1, \dots, a_n$  be formal variables, where we shorthand  $\mathbf{a} = (a_1, \dots, a_n)$ . We use the standard notations  $\mathbb{F}[\mathbf{a}, x]$  for the ring of polynomials over  $\mathbb{F}$  in the variables  $\mathbf{a}, x$ ;  $\mathbb{F}(\mathbf{a})$  for the field of rational functions over  $\mathbb{F}[\mathbf{a}]$ ; and  $\mathbb{F}(\mathbf{a})[x]$  for the ring of univariate polynomials in  $x$  over  $\mathbb{F}(\mathbf{a})$ . Given a set  $S \subset [n]$  define a polynomial  $p = p(S) \in \mathbb{F}[\mathbf{a}, x]$  as follows:

$$p(\mathbf{a}, x) := \prod_{i \in S} (x - a_i).$$

Given a set system  $\mathcal{S} = \{S_1, \dots, S_k\}$  define  $P(\mathcal{S}) := \{p(S_1), \dots, p(S_k)\}$ .

Let  $\mathcal{S} = \{S_1, \dots, S_k\}$  be a set system which satisfies (1). It is possible to assume without loss of generality that each  $S_i$  is maximal, namely that  $|S_i| = k - 1$  for all  $i \in [k]$ . For example, if we are allowed to increase  $n$  then we can replace each  $S_i$  with  $S_i \cup T_i$  where  $|T_i| = k - 1 - |S_i|$  and  $T_1, \dots, T_k, [n]$  are pairwise disjoint. An improved reduction is given in [DSY14] which does not require increasing  $n$ .

Either way, under this assumption the polynomials  $P(\mathcal{S})$  form a set of  $k$  polynomials of degree  $k - 1$ , which we denote by  $p_1, \dots, p_k$ . Define the  $k \times n$  matrix  $G$  as  $G_{i,j} = p_i(a_j)$ . Note that entries of  $G$  are polynomials in  $\mathbb{F}[\mathbf{a}]$ . The condition that all  $k \times k$  minors of  $G$  are nonsingular is equivalent to the condition that the polynomials  $P(\mathcal{S})$  are linearly independent over  $\mathbb{F}(\mathbf{a})$  (here, we view the polynomials as elements of  $\mathbb{F}(\mathbf{a})[x]$  instead of as elements of  $\mathbb{F}[\mathbf{a}, x]$ ). If this is the case, then one can use the Schwartz-Zippel lemma and show that the formal variables  $a_1, \dots, a_n$  can be replaced with distinct field elements from  $\mathbb{F}$ , while still maintaining the property that all  $k \times k$  minors of  $G$  are nonsingular. The bound on the field size  $|\mathbb{F}| \geq n + k - 1$  arises from the degrees of the polynomials obtained in the process. For details we refer to the original paper [DSY14].

This motivated [DSY14] to propose the following algebraic conjecture, which implies Conjecture 1.1.

**Conjecture 1.2** (Algebraic GM-MDS conjecture [DSY14]). *Let  $S_1, \dots, S_k \subset [n]$  be a set system which satisfies the MDS condition, and where  $|S_i| = k - 1$  for all  $i$ . Then the set of polynomials  $P(\mathcal{S})$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ .*

We remark that given any polynomials  $p_1, \dots, p_k \in \mathbb{F}[\mathbf{a}, x]$  (for example, the polynomials appearing in  $P(\mathcal{S})$ ), an equivalent condition to the polynomials being linearly independent over  $\mathbb{F}(\mathbf{a})$  is the following: for any polynomials  $w_1, \dots, w_k \in \mathbb{F}[\mathbf{a}]$ , not all zero, it holds that

$$\sum_{i=1}^k w_i(\mathbf{a}) p_i(\mathbf{a}, x) \neq 0.$$

Following [DSY14], several works [HHYD14, HS17, YH18] attempted to resolve the GM-MDS conjecture. They showed that Conjecture 1.2 holds in several special cases, but the general case remained open. In this work we prove Conjecture 1.2, which implies Conjecture 1.1.

## 1.2 A generalized conjecture

We start by considering a more general condition. Let  $v \in \mathbb{N}^n$  be a vector, where  $\mathbb{N} = \{0, 1, 2, \dots\}$  stands for non-negative integers. The coordinates of  $v$  are  $v = (v(1), \dots, v(n))$ . We shorthand  $|v| = \sum v(i)$ . Given vectors  $v_1, \dots, v_m \in \mathbb{N}^n$  define  $\bigwedge v_i \in \mathbb{N}^n$  to be their coordinate-wise minimum:

$$\bigwedge_{i \in [m]} v_i := (\min(v_1(1), \dots, v_m(1)), \dots, \min(v_1(n), \dots, v_m(n))).$$

Note that if  $v_1, \dots, v_m \in \{0, 1\}^n$  are indicator vectors of sets  $S_1, \dots, S_m \subset [n]$ , then  $\bigwedge v_i$  is the indicator vector of  $\cap S_i$ .

Given a parameter  $k > |v|$  define a set of polynomials in  $\mathbb{F}[\mathbf{a}, x]$ :

$$P(k, v) := \left\{ \prod_{j \in [n]} (x - a_j)^{v(j)} x^e : e = 0, \dots, k - 1 - |v| \right\}.$$

Note that  $P(k, v)$  consists of  $k - |v|$  polynomials of degree  $\leq k - 1$ , which form a basis for the linear space of polynomials of degree  $\leq k - 1$  which have  $v(j)$  roots at each  $a_j$ . Furthermore, note that if  $v$  is the indicator vector of a set  $S \subset [n]$  of size  $|S| = k - 1$ , then  $P(k, v) = \{p(S)\}$ . Given a set of vectors  $\mathcal{V} = \{v_1, \dots, v_m\} \subset \mathbb{N}^n$  define

$$P(k, \mathcal{V}) := P(k, v_1) \cup \dots \cup P(k, v_m).$$

We use in this paper the convention that set union can result in a multiset. So for example, if the same polynomial appears in multiple  $P(k, v_i)$  then it appears multiple times in  $P(k, \mathcal{V})$ . Under this assumption we always have the identity:

$$|P(k, \mathcal{V})| = |P(k, v_1)| + \dots + |P(k, v_m)|.$$

The following definition is the natural extension of the MDS condition (1) to vectors.

**Definition 1.3** (Property  $V(k)$ ). *Let  $\mathcal{V} = \{v_1, \dots, v_m\} \subset \mathbb{N}^n$  and  $k \geq 1$  be an integer. We say that  $\mathcal{V}$  satisfies  $V(k)$  if it satisfies:*

- (i)  $|v_i| \leq k - 1$  for all  $i \in [m]$ .
- (ii) For all  $I \subseteq [m]$ ,  $\sum_{i \in I} (k - |v_i|) + |\bigwedge_{i \in I} v_i| \leq k$ .

Note that when  $m = k$  and  $v_1, \dots, v_k$  are indicators of sets  $S_1, \dots, S_k \subset [n]$  of size  $|S_i| = k - 1$ , then property  $V(k)$  is equivalent to the MDS condition for  $S_1, \dots, S_k$ .

Observe that in general, if  $\mathcal{V}$  satisfies  $V(k)$  then  $P(k, \mathcal{V})$  contains  $\sum_{i=1}^m (k - |v_i|) \leq k$  polynomials of degree  $\leq k - 1$ . The following conjecture is the natural extension of Conjecture 1.2 to vectors.

**Conjecture 1.4.** *Let  $\mathcal{V} \subset \mathbb{N}^n$  and  $k \geq 1$ . Assume that  $\mathcal{V}$  satisfies  $V(k)$ . Then the polynomials in  $P(k, \mathcal{V})$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ .*

A clarifying remark: as we view the set  $P(k, \mathcal{V})$  as a multiset, Conjecture 1.4 (and Theorem 1.6 below) imply in particular that the polynomials in  $P(k, \mathcal{V})$  are all distinct, so  $P(k, \mathcal{V})$  is in fact a set.

### 1.3 An intermediate case

We prove Conjecture 1.4 under an additional assumption, which is sufficient to prove Conjecture 1.1. It is still open to prove Conjecture 1.4 in full generality.

**Definition 1.5** (Property  $V^*(k)$ ). *Let  $\mathcal{V} = \{v_1, \dots, v_m\} \subset \mathbb{N}^n$  and  $k \geq 1$  be an integer. We say that  $\mathcal{V}$  satisfies  $V^*(k)$  if it satisfies  $V(k)$ , and additionally it satisfies:*

(iii)  $v_i \in \{0, 1\}^{n-1} \times \mathbb{N}$  for all  $i \in [m]$ . Namely, all coordinates in  $v_i$ , except perhaps the last, are in  $\{0, 1\}$ .

**Theorem 1.6.** *Let  $\mathcal{V} \subset \mathbb{N}^n$  and  $k \geq 1$ . Assume that  $\mathcal{V}$  satisfies  $V^*(k)$ . Then the polynomials  $P(k, \mathcal{V})$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ .*

Conjecture 1.2 follows directly from Theorem 1.6. If  $S_1, \dots, S_k \subset [n]$  are sets which satisfy the assumptions of Conjecture 1.2, then their indicator vectors  $v_1, \dots, v_k \in \{0, 1\}^n$  satisfy the assumptions of Theorem 1.6, and hence  $P(\{S_1, \dots, S_k\}) = P(k, \{v_1, \dots, v_k\})$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ .

## 2 Proof of Theorem 1.6

Let  $n, k \geq 1$ . Let  $\mathcal{V} = \{v_1, \dots, v_m\} \subset \mathbb{N}^n$  which satisfies  $V^*(k)$ . We will prove that the polynomials  $P(k, \mathcal{V})$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ .

To that end, we assume that  $\mathcal{V}$  is a minimal counter-example and derive a contradiction. Concretely, the underlying parameters are  $n, k, m$  and  $d = |P(k, \mathcal{V})| = \sum k - |v_i|$ . We will assume that if  $\mathcal{V}'$  is a set of vectors with corresponding parameters  $n' \leq n, k' \leq k, m' \leq m, d' \leq d$  with at least one of the inequalities being sharp, then Theorem 1.6 holds for  $\mathcal{V}'$ . In particular, we assume that  $m \geq 2$ , as Theorem 1.6 clearly holds when  $m = 1$ .

To help the reader, we note that the following lemmas construct such  $\mathcal{V}'$  with the following parameters:

- Lemma 2.2:  $n, k - 1, m, d$ .

- Lemma 2.4:  $n, k, e, d'$  and  $n, k, m - e + 1, d''$  with  $2 \leq e \leq m - 1$  and  $d', d'' < d$ .
- Lemma 2.5:  $n - 1, k, m, d$ .
- Lemma 2.6:  $n, k, m, d - 1$ .

We use the following notation to simplify the presentation:

$$v_I := \bigwedge_{i \in I} v_i \quad I \subseteq [m].$$

We introduce sometimes in the proofs an auxiliary set  $\mathcal{V}' = \{v'_1, \dots, v'_{m'}\}$ , in which case  $v'_I$  for  $I \subseteq [m']$  are defined analogously. Below, when we say that  $\mathcal{V}$  or  $\mathcal{V}'$  satisfy (i), (ii) or (iii), we mean the relevant items in the definition of  $V^*(k)$ .

Given two vectors  $u, v \in \mathbb{N}^n$  we denote  $u \leq v$  if  $u(i) \leq v(i)$  for all  $i \in [n]$ .

**Lemma 2.1.** *There do not exist distinct  $i, j \in [m]$  such that  $v_i \leq v_j$ .*

*Proof.* Assume the contrary. Applying (i) to  $j$  gives  $|v_j| \leq k - 1$ . Applying (ii) to  $I = \{i, j\}$  gives

$$(k - |v_i|) + (k - |v_j|) + |v_i \wedge v_j| \leq k.$$

As  $v_i \leq v_j$  we have  $v_i \wedge v_j = v_i$ , and hence obtain that  $k - |v_j| \leq 0$ , a contradiction.  $\square$

Lemma 2.1 implies in particular that  $n \geq 2$ . This is since if  $n = 1$  then necessarily  $m = 1$ , as otherwise there would be  $i, j$  for which  $v_i \leq v_j$ . So we assume  $n \geq 2$  from now on.

**Lemma 2.2.**  $\bigwedge_{i \in [m]} v_i = 0$ .

*Proof.* Assume not. Then there exists a coordinate  $j \in [n]$  with  $v_i(j) \geq 1$  for all  $i \in [m]$ . Define a new set of vectors  $\mathcal{V}' = \{v'_1, \dots, v'_m\} \subset \mathbb{N}^n$  as follows:

$$v'_i := (v_i(1), \dots, v_i(j-1), v_i(j) - 1, v_i(j+1), \dots, v_i(n)).$$

In words,  $v'_i$  is defined from  $v_i$  by decreasing coordinate  $j$  by 1.

We first show that  $\mathcal{V}'$  satisfies  $V^*(k-1)$ . Note that  $|v'_i| = |v_i| - 1$ . It clearly satisfies (i),(iii). To show that it satisfies (ii) let  $I \subseteq [m]$ . We have

$$\sum_{i \in I} (k - 1 - |v'_i|) + |v'_I| = \sum_{i \in I} (k - |v_i|) + |v_I| - 1 \leq k - 1.$$

As we showed that  $\mathcal{V}'$  satisfies  $V^*(k-1)$ , the minimality of  $\mathcal{V}$  implies that the polynomials  $P(k-1, \mathcal{V}')$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ . The lemma follows as it is simple to verify that

$$P(k, \mathcal{V}) = \{p(\mathbf{a}, x)(x - a_j) : p \in P(k-1, \mathcal{V}')\}.$$

In particular, the linear independence of  $P(k-1, \mathcal{V}')$  implies the linear independence of  $P(k, \mathcal{V})$ .  $\square$

**Definition 2.3** (Tight constraint). A set  $I \subseteq [m]$  is tight for  $\mathcal{V}$  if property (ii) holds with equality for  $I$ . Namely if

$$\sum_{i \in I} (k - |v_i|) + |v_I| = k.$$

Note that if  $|I| = 1$  then  $I$  is always a tight constraint. The following lemma is an extension of Lemma 2(i) in [YH18]. It shows that in a minimal counter-example there are no tight sets, except for singletons and perhaps the whole set.

**Lemma 2.4.** *If  $I \subseteq [m]$  is a tight constraint, then  $|I| = 1$  or  $|I| = m$ .*

*Proof.* Assume towards a contradiction that there exist a tight  $I$  with  $1 < |I| < m$ . We will use the minimality of  $\mathcal{V}$  to derive a contradiction. Assume for simplicity of notation that  $I = \{e, \dots, m\}$  for  $2 \leq e \leq m - 1$ . Define a new set of vectors  $\mathcal{V}' = \{v'_1, \dots, v'_e\}$  given by

$$v'_1 := v_1, \dots, v'_{e-1} := v_{e-1}, v'_e := v_e.$$

We first show that  $\mathcal{V}'$  satisfies  $V^*(k)$ . It clearly satisfies (i) and (iii). To see that it satisfies (ii) let  $I' \subseteq [e]$ . If  $e \notin I'$  then  $\mathcal{V}'$  satisfies (ii) for  $I'$  as it is same condition as for  $\mathcal{V}$ , so assume  $e \in I'$ . Let  $I'' = I' \cup \{e + 1, \dots, m\}$ . Then

$$\sum_{i \in I'} (k - |v'_i|) + |v'_{I'}| = \sum_{i \in I''} (k - |v_i|) + |v_{I''}| \leq k,$$

where the equality holds since  $k - |v'_e| = \sum_{i \in I} (k - |v_i|)$  since we assume  $I$  is tight, and since by definition of  $I''$  we have  $v'_I = v_{I''}$ .

As we assume that  $\mathcal{V}$  is a minimal counter-example for Theorem 1.6, the theorem holds for  $\mathcal{V}'$ . So, the polynomials  $P(k, \mathcal{V}')$  are linearly independent. Observe that  $|P(k, \mathcal{V}')| = |P(k, \mathcal{V})|$  since

$$|P(k, \mathcal{V}')| = \sum_{i \in [e]} (k - |v'_i|) = \sum_{i \in [m]} (k - |v_i|) = |P(k, \mathcal{V})|.$$

Thus, it will suffice to prove that  $P(k, \mathcal{V})$  and  $P(k, \mathcal{V}')$  span the same space of polynomials over  $\mathbb{F}(\mathbf{a})$ . To that end, it suffices to prove that  $F := P(k, \{v_e, \dots, v_m\})$  and  $F' := P(k, v'_e)$  span the same space of polynomials.

Let us shorthand  $v = v'_e$ . Define the polynomial  $p(\mathbf{a}, x) := \prod_{j \in [n]} (x - a_j)^{v(j)}$ . Observe that  $p$  divides all polynomials in  $F, F'$ . Moreover,  $F' = \{p(\mathbf{a}, x)x^d : d = 0, \dots, k - 1 - |v|\}$  spans the linear space of all multiples of  $p$  of degree  $\leq k - 1$ . As  $|F| = |F'|$  it suffices to prove that  $F$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ , as then they must span the same linear space. However, this follows from the minimality of  $\mathcal{V}$ , since  $F = P(k, \mathcal{V}'')$  for  $\mathcal{V}'' = \{v_e, \dots, v_m\}$ .  $\square$

The following lemma identifies a concrete vector that must exist in a minimal counter-example. It is in its proof that we actually use the assumption that  $\mathcal{V}$  satisfies (iii), namely  $V^*(k)$  and not merely  $V(k)$ .

**Lemma 2.5.** *There exists  $i \in [m]$  such that  $v_i = (1, \dots, 1, 0)$ .*

*Proof.* Lemma 2.2 guarantees that there exists  $i^* \in [m]$  for which  $v_{i^*}(n) = 0$ . We will prove that  $v_{i^*} = (1, \dots, 1, 0)$ . If not, then by (iii) there exists  $j^* \in [n-1]$  be such that  $v_{i^*}(j^*) = 0$ . For simplicity of notation assume that  $i^* = m, j^* = n-1$ . Define a new set of vectors  $\mathcal{V}' = \{v'_1, \dots, v'_m\} \subset \mathbb{N}^{n-1}$  as follows:

$$v'_i := (v_i(1), \dots, v_i(n-2), v_i(n-1) + v_i(n)).$$

In words,  $v'_i \in \mathbb{N}^{n-1}$  is obtained by adding the last two coordinates of  $v_i \in \mathbb{N}^n$ .

We first show that  $\mathcal{V}'$  satisfies  $V^*(k)$ . Note that  $|v'_i| = |v_i|$ . It clearly satisfies (i),(iii). To show that it satisfies (ii) let  $I \subseteq [m]$ . Note that (ii) always holds if  $|I| = 1$ , so we may assume  $|I| > 1$ . We have by definition

$$\sum_{i \in I} (k - |v'_i|) + |v'_I| - v'_I(n-1) = \sum_{i \in I} (k - |v_i|) + |v_I| - v_I(n-1) - v_I(n). \quad (2)$$

First, consider first the case where  $|I| < m$ . Lemma 2.4 gives that  $I$  is not tight, and hence

$$\sum_{i \in I} (k - |v_i|) + |v_I| \leq k - 1.$$

As  $\mathcal{V}$  satisfies (iii) we have  $v_i(n-1) \in \{0, 1\}$  for all  $i$ . This implies  $v_I(n-1) \in \{0, 1\}$  and  $v'_I(n-1) \in \{v_I(n), v_I(n) + 1\}$ . So Equation (2) gives

$$\sum_{i \in I} (k - |v'_i|) + |v'_I| \leq \sum_{i \in I} (k - |v_i|) + |v_I| + 1 \leq k.$$

Next, consider the case of  $|I| = m$ . As  $v_m(n-1) = v_m(n) = 0$  we have  $v'_m(n-1) = 0$  and hence  $v_I(n-1) = v_I(n) = v'_I(n-1) = 0$ . Equation (2) then gives

$$\sum_{i \in I} (k - |v'_i|) + |v'_I| = \sum_{i \in I} (k - |v_i|) + |v_I| \leq k.$$

As we showed that  $\mathcal{V}'$  satisfies  $V^*(k)$ , the minimality of  $\mathcal{V}$  implies that the polynomials  $P(k, \mathcal{V}')$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ . We next show that this implies that  $P(k, \mathcal{V})$  are also linearly independent over  $\mathbb{F}(\mathbf{a})$ .

Let  $s_i := k - |v_i|$  for  $i \in [m]$ . We have  $P(k, \mathcal{V}) = \{p_{i,e} : i \in [m], e \in [s_i]\}$  and  $P(k, \mathcal{V}') = \{p'_{i,e} : i \in [m], e \in [s_i]\}$  where

$$p_{i,e}(\mathbf{a}, x) := x^{e-1} \prod_{j \in [n-2]} (x - a_j)^{v_i(j)} \cdot (x - a_{n-1})^{v_i(n-1)} (x - a_n)^{v_i(n)},$$

$$p'_{i,e}(\mathbf{a}, x) := x^{e-1} \prod_{j \in [n-2]} (x - a_j)^{v_i(j)} \cdot (x - a_{n-1})^{v_i(n-1) + v_i(n)}.$$

Observe that  $p'_{i,e}$  can be obtained from  $p_{i,e}$  by substituting  $a_{n-1}$  for  $a_n$ . Namely

$$p'_{i,e}(a_1, \dots, a_{n-1}, x) = p_{i,e}(a_1, \dots, a_{n-1}, a_{n-1}, x).$$

Assume towards a contradiction that  $\{p_{i,e}\}$  are linearly dependent over  $\mathbb{F}(\mathbf{a})$ . Equivalently, there exist polynomials  $w_{i,e}(\mathbf{a})$ , not all zero, such that

$$\sum_{i \in [m]} \sum_{j \in [s_i]} w_{i,e}(\mathbf{a}) p_{i,e}(\mathbf{a}, x) = 0.$$

We may assume that the polynomials  $\{w_{i,e}\}$  do not all have a common factor, as otherwise we can divide them by it. Let  $w'_{i,e}(\mathbf{a})$  be obtained from  $w_{i,e}(\mathbf{a})$  by substituting  $a_{n-1}$  for  $a_n$ . That is,  $w'_{i,e}(a_1, \dots, a_{n-1}) = w_{i,e}(a_1, \dots, a_{n-1}, a_{n-1})$ . Then we obtain

$$\sum_{i \in [m]} \sum_{j \in [s_i]} w'_{i,e}(\mathbf{a}) p'_{i,e}(\mathbf{a}, x) = 0.$$

As the polynomials  $\{p'_{i,e}\}$  are linearly independent over  $\mathbb{F}(\mathbf{a})$ , this implies that  $w'_{i,e} \equiv 0$  for all  $i, e$ . That is, the polynomials  $w_{i,e}$  satisfy

$$w_{i,e}(a_1, \dots, a_{n-1}, a_{n-1}) \equiv 0.$$

This implies that  $(a_{n-1} - a_n)$  divides  $w_{i,e}$  for all  $i, e$ , which is a contradiction to the assumption that  $\{w_{i,e}\}$  do not all have a common factor.  $\square$

Lemma 2.5 implies that the vector  $(1, \dots, 1, 0)$  belongs to  $\mathcal{V}$ . Without loss of generality, we may assume that it is  $v_m$ . This implies that  $v_i(n) \geq 1$  for all  $i \in [m-1]$ , as otherwise we would have  $v_i \leq v_m$ , violating Lemma 2.1.

**Lemma 2.6.**  $n = k$ .

*Proof.* Let  $v_m = (1, \dots, 1, 0)$ . We know by (i) that  $n - 1 = |v_m| \leq k - 1$ , so  $n \leq k$ . Assume towards a contradiction that  $n < k$ . Define a new set of vectors  $\mathcal{V}' = \{v'_1, \dots, v'_m\} \subset \mathbb{N}^n$  as follows:

$$v'_1 := v_1, \dots, v'_{m-1} := v_{m-1}, v'_m := (1, \dots, 1, 1).$$

In words, we increase the last coordinate of  $v_m$  by 1.

We claim that  $\mathcal{V}'$  satisfies  $V^*(k)$ . It satisfies (i) by our assumption that  $|v'_m| = n \leq k - 1$ , and it satisfies (iii) clearly. To show that it satisfies (ii), let  $I \subseteq [m]$ . If  $m \notin I$  then it clearly satisfies (ii) for  $I$ , as it is the same constraint as for  $\mathcal{V}$ , so assume  $m \in I$ . In this case we have

$$\sum_{i \in I} (k - |v'_i|) + |v'_I| = \left( \sum_{i \in I} (k - |v_i|) - 1 \right) + (|v_I| + 1) \leq k.$$

Note that  $|P(k, \mathcal{V}')| = |P(k, \mathcal{V})| - 1$ . As  $\mathcal{V}$  is a minimal counter-example, we have that  $\mathcal{V}'$  satisfies  $V^*(k)$ . Let  $p(\mathbf{a}, x) := \prod_{j \in [n-1]} (x - a_j)$ . The construction of  $\mathcal{V}'$  satisfies that

$$P(k, \mathcal{V}) = P(k, \mathcal{V}') \cup \{p\}.$$

Denote for simplicity of presentation the polynomials of  $P(k, \mathcal{V}')$  by  $p_1, \dots, p_{d-1}$ , where  $d = |P(k, \mathcal{V})|$ . Assume that the polynomials  $P(k, \mathcal{V})$  are linearly dependent. As  $P(k, \mathcal{V})$  are



linearly independent, it implies that there exist polynomials  $w, w_1, \dots, w_{d-1} \in \mathbb{F}[\mathbf{a}]$ , where  $w \neq 0$ , such that

$$w(\mathbf{a})p(\mathbf{a}, x) + \sum_{i=1}^{d-1} w_i(\mathbf{a})p_i(\mathbf{a}, x) \equiv 0.$$

Note that by construction,  $v'_i(n) \geq 1$  for all  $i \in [m]$ . This implies that  $p_1, \dots, p_{d-1}$  are all divisible by  $(x - a_n)$ , while  $p$  is not. Substituting  $x = a_n$  then gives  $w \equiv 0$ , a contradiction.  $\square$

We can now reach a contradiction to  $\mathcal{V}$  being a counter-example. We know that  $v_m = (1, \dots, 1, 0)$  with  $|v_m| = n - 1 = k - 1$ . Let  $\mathcal{V}' = \{v_1, \dots, v_{m-1}\}$ . As it satisfies  $V^*(k)$  we have that the polynomials  $P(k, \mathcal{V}')$  are linearly independent. Moreover, as  $|v_m| = k - 1$  we have  $P(k, v_m) = \{p\}$  where  $p(\mathbf{a}, x) = \prod_{j \in [n-1]} (x - a_j)$ . Note that all polynomials in  $P(k, \mathcal{V}')$  are divisible by  $(x - a_n)$ , while  $p$  is not. So by the same argument as in the proof of Lemma 2.6,  $P(k, v_m)$  cannot be linearly dependent of  $P(k, \mathcal{V}')$ . So  $P(k, \mathcal{V})$  are linearly independent.

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