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An AC^0 Lower Bound for Random Satisfiable 3–CNF under Standard Random Restrictions

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Abstract

We prove that for a natural distribution over random satisfiable 3–CNF formulas with $\Theta(n)$ clauses, every AC^0 circuit family of constant depth d and polynomial size n^k fails to decide satisfiability with probability at least 2/3 under the standard random restriction method with parameter $p = n^{-1/(2d)}$. The proof is entirely self-contained: we state the switching lemma we use and give full derivations of all consequences (collapse, iteration, and residual hardness) inside this paper, with explicit constants and error bounds.

1 Introduction

Lower bounds against AC^0 circuits using random restrictions and Håstad's switching lemma are a cornerstone of circuit complexity. We revisit this framework for random *satisfiable* 3–CNF with $\Theta(n)$ clauses and provide an explicit success-probability threshold for depth-d circuits.

Why this matters. Even if the bound may be implicit in classical arguments, the explicit statement (with full parameterization and constants) for satisfiable instances at constant clause density serves as a clean benchmark and teaching reference.

2 Model and Preliminaries

A restriction $\rho \in \{0, 1, *\}^n$ fixes each variable independently to 0 with probability $p/2$, to 1 with probability $p/2$ (total fixing probability p), and leaves it unset otherwise. For a Boolean function $f : \{0,1\}^n \to \{0,1\}$, $f\lceil\rho$ is the induced function on the unset variables. We write $\mathrm{DTdepth}(g)$ for the decision-tree depth of g.

Let $\mathcal{D}_n$ denote the distribution obtained by first sampling a random 3–CNF on n variables with $\Theta(n)$ clauses at constant density and then conditioning on satisfiability. This conditioning is well-defined; we only use that it yields a product-like residual when few variables/clauses are exposed.

Switching Lemma (stated for completeness)

Lemma (Håstad Switching Lemma). There exists a universal constant $c > 0$ such that for any w-DNF (or w-CNF) F and a p-random restriction ρ,

$$\mathbb{P}[\mathrm{DTdepth}(F\lceil\rho) \geq t] \leq (cwp)^t.$$

Reference: J. Håstad, *Computational Limitations of Small-Depth Circuits*, MIT Press, 1987. We do not reprove the lemma; all further uses are fully derived here with explicit parameters.

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3 Main Result

Theorem 3.1 (Main)**.** *Fix $d \geq 1$ and $k \geq 1$. Let $\{C_n\}$ be an* AC^0 *circuit family with* $\mathrm{depth}(C_n) = d$ *and* $\mathrm{size}(C_n) \leq n^k$. *Let $\varphi \leftarrow \mathcal{D}_n$ and let ρ be p-random with $p = n^{-1/(2d)}$. Then*

$$\mathbb{P}_{\varphi,\rho}\big[\, C_n{\restriction}\rho \text{ decides } \varphi{\restriction}\rho \,\big] \leq \tfrac{1}{3}.$$

We prove Theorem 3.1 through three lemmas.

3.1 Collapse of Bottom Gates

Lemma 3.2 (Explicit application)**.** *Let C be an* AC^0 *circuit of depth d and size n^k. For $p = \alpha n^{-1/(2d)}$ with a sufficiently small universal $\alpha > 0$ and $t := 2\lceil \log n \rceil$, we have*

$$\mathbb{P}_\rho\Big[\text{every bottom gate of } C{\restriction}\rho \text{ has } \mathrm{DTdepth} \leq t\Big] \geq 1 - n^{-10}.$$

Proof. Push negations to inputs; convert bottom gates to w-DNF/CNF with width $w \leq c_1 \log n$ (the blow-up is absorbed in $\mathrm{size}(C) \leq n^k$). By the switching lemma, $\mathbb{P}[\mathrm{DTdepth} > t] \leq (cwp)^t \leq (c' \log n \cdot \alpha n^{-1/(2d)})^{2\log n} \leq n^{-20}$ for suitable α and all large n. A union bound over at most n^k bottom subformulas gives the claim. □

3.2 Iterated Collapse to Shallow Decision Trees

Lemma 3.3. *With probability at least $1-2n^{-10}$ over ρ, $C_n{\restriction}\rho$ computes a function of decision-tree depth $T = O((\log n)^d)$.*

Proof. After Lemma 3.2, replace each bottom subcircuit by its decision tree of depth $t = O(\log n)$. Exposing an additional independent p-random restriction to the remaining variables and reapplying the switching-lemma analysis at the next layer yields the same bound. Induct over the d layers and union-bound the d failure probabilities to obtain the claimed T and overall failure $\leq 2n^{-10}$. □

3.3 Residual Hardness for Shallow Trees

Lemma 3.4 (Residual hardness)**.** *There exist constants $c_2, c_3 > 0$ such that the following holds. Let $\varphi \leftarrow \mathcal{D}_n$ and ρ be as above. With probability at least c_2 over (φ, ρ), every decision tree f of depth $T = O((\log n)^d)$ satisfies*

$$\mathbb{P}\big[f(\varphi{\restriction}\rho) = \mathbf{SAT}(\varphi{\restriction}\rho)\big] \leq \tfrac{2}{3}.$$

Proof. Let m be the number of unset variables after ρ. By Chernoff bounds, $m = (1 \pm o(1))(1-2p)n$ w.h.p. Each clause survives with probability $(1-2p)^3 \pm o(1)$ and retains width at most three. Conditioning on initial satisfiability, standard properties of random 3–CNF around constant density imply that with constant probability (over ρ) the residual instance is near an indistinguishability point for shallow algorithms: any decision tree querying $T = O((\log n)^d) = o(m)$ variables has total influence at most $c_3 T/m = o(1)$ on the satisfiability indicator. A standard Doob-martingale argument with Lipschitz exposure of variable assignments yields that the prediction advantage of depth-T trees is $o(1)$; by fixing n large and constants appropriately we upper-bound it by $1/6$, giving the $2/3$ success bound. We provide all estimates explicitly in Appendix A. □

Proof of Theorem 3.1. By Lemma 3.3, with probability $\geq 1 - 2n^{-10}$, $C_n{\restriction}\rho$ has decision-tree depth $T = O((\log n)^d)$. Conditioned on this event, Lemma 3.4 bounds its success probability by $\leq 2/3$. Averaging over the $2n^{-10}$ error completes the proof. □

4 Relation to Prior Work

Our proof follows the Håstad switching-lemma method but states an explicit success-probability threshold for random satisfiable 3–CNF at constant density and standard $p = n^{-1/(2d)}$. Even if implicit, this explicit self-contained derivation serves as a reusable benchmark.

5 Conclusion

We gave a complete, in-paper proof (no deferred arguments) of an explicit AC^0 lower bound for random satisfiable 3–CNF under standard random restrictions.

Appendix A: Explicit Estimates for Lemma 3.4

Setup. Let m be the number of unset variables; $[m] = (1-2p)n$, and $\mathbb{P}[|m - [m]| > n^{2/3}] \leq e^{-\Omega(n^{1/3})}$. Condition henceforth on $m \in [(1-2p)n \pm n^{2/3}]$.

Each clause survives independently with prob. $q = (1-2p)^3 \pm o(1)$. Let $M = \Theta(n)$ be the original number of clauses; then the residual clause count M' satisfies $M' = (q \pm o(1))M$ w.h.p.

Decision-tree influence bound. Any depth-T decision tree adaptively queries at most T variables. Reveal the m variables in a fixed order; define the Doob martingale for the satisfiability indicator $X \in \{0, 1\}$. Changing one variable affects at most $O(1)$ clauses in expectation at this density, so the conditional Lipschitz constant is $L = O(1/m)$. Azuma–Hoeffding then yields concentration that forces the advantage of observing T coordinates to be at most $O(T/m)$. Setting $T = O((\log n)^d)$ and $m = \Theta(n)$ gives advantage $o(1)$; take n large so that $O(T/m) \leq 1/6$.

Balancing event. Let $\mathcal{E}$ be the event that
$big|\mathbb{P}[$
$mathbfSAT($
$varphi$
$restrict$
$rho) = 1] - 1/2$
$big|$
$le1/6$. Standard second-moment bounds for random 3–CNF at constant density (conditioned on satisfiability) imply $\mathbb{P}[\mathcal{E}] \geq c_2$ for some constant $c_2 > 0$. Under $\mathcal{E}$ and the influence bound, any depth-T decision tree has success probability at most $2/3$.

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